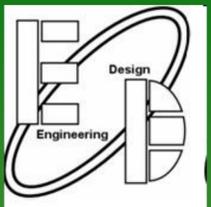




Volume Constrained Polyhedronizations of Point Sets in 3-Space



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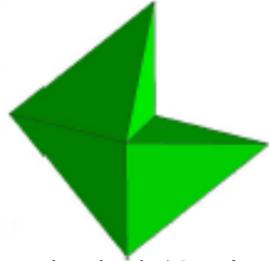
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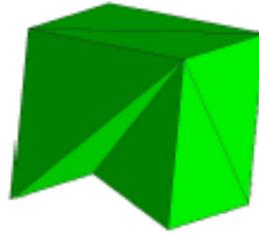
Minimum Volume Polyhedronization of Platonic Point sets



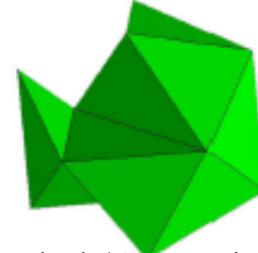
Tetrahedral (Optimal)



Octahedral (Optimal)



Cube (Optimal)



Icosahedral (Approximate)



Dodecahedral (Approximate)

Polyhedronization

•Definition

“Given a finite set of points in R^3 , polyhedronization deals with constructing a simple polyhedron such that the vertices of the polyhedron are precisely the given points.”

•Applications

- Molecular polyhedron structure synthesis.
- Boundary representation of input points in Computer Graphics, Computer Vision & Distance Image Processing.

Algorithm

•RAA_MINVP-Randomized Approximation Algorithm

Let $S = \{p_0, p_1, \dots, p_{(n-1)}\}$ denotes the point set.

Initialization

Select four points uniformly at random from S and form an initial tetrahedron P .

Iterations

In each iteration, it chooses one point q uniformly at random from $S \setminus P$. Determines the position of q relative to the previous polyhedron P and does one of the following.

1. q lies interior to P ? \rightarrow exclude from P , the largest volume tetrahedron that q makes with any of the visible faces of P .
2. q lies exterior to P ? \rightarrow add to P , the smallest volume tetrahedron that q forms with any of the visible faces of P .
3. q lies on an edge of P ? \rightarrow divide the adjacent faces of that edge into four new faces by including q as the common vertex of all the four faces.
4. q lies on a face of P ? \rightarrow divide the face into three new faces by including q as the common vertex of all the three faces.

Termination

Once the iterations are completed, algorithm returns the final polyhedron (The set of faces).

•RAA_MAXVP

The initial polyhedron is the convex hull of S . The iterations are pretty much similar to the iterations of RAA_MINVP. Both differs only in steps 1 & 2.

- q lies interior to P ? \rightarrow exclude from P , the smallest volume tetrahedron that q makes with any of the visible faces of P and vice versa.

Problem Statement

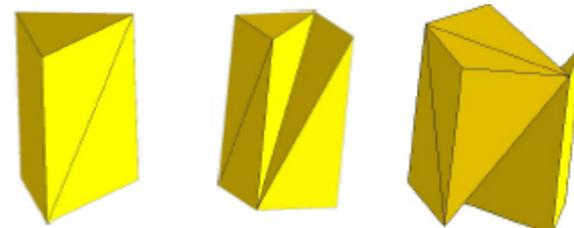
•FACE problem by S.P Fekete [FP93]

“Let $2 \leq d$ and $1 \leq k \leq d$. Given a finite set S of points in d -dimensional Euclidean space. Among all simple polyhedra that are feasible for vertex set S , find one with the smallest volume of its k -dimensional faces.”

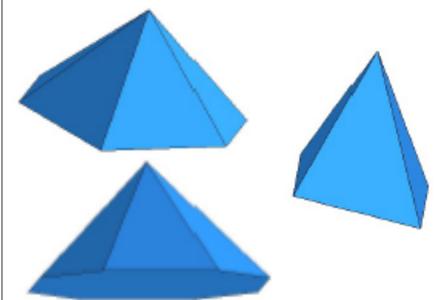
•Minimal(Maximal) Volume Polyhedronization (MINVP (MAXVP))

“Given a finite set S of n points in R^3 , find the simple polyhedron with the smallest (largest) volume from all the simple polyhedra (having triangular faces) that are feasible for the vertex set S .”

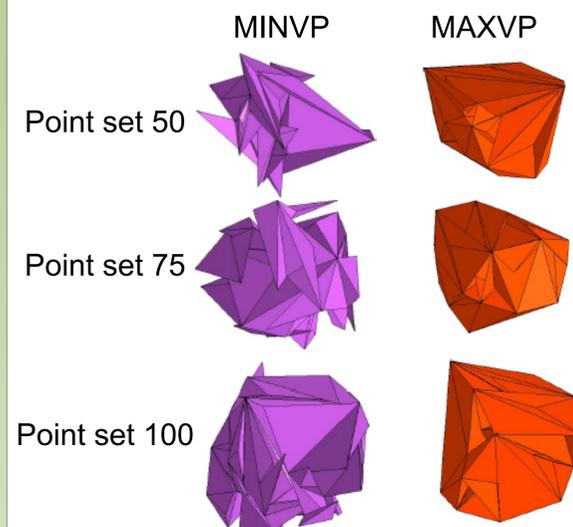
Results



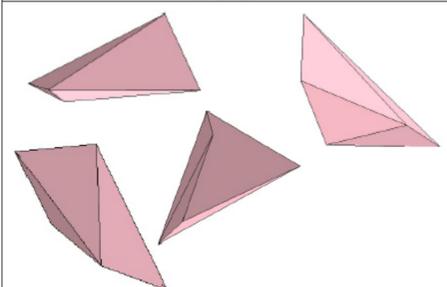
Approximate MINVPs generated for **Prismatic** Point Sets.



MINVPs and/or MAXVPs generated for **Pyramid** Point Sets.



Approximate MINVPs & MAXVPs generated for Point Sets of different sizes.



Optimal MINVPs generated by RAA_MINVP algorithm for point sets of size 5. The results are verified using brute force approach.

Future Work

•To address the following questions:

- What are the performance guarantees of both the algorithms?
- Does there exist an input configuration for which the approach fails for every possible ordering of points?

References

[FP93] FEKETE S. P., PULLEYBLANK W. R.: Area optimization of simple polygons. In *Proc. 9th Annu. ACM Sympos. Computational Geometry. (1993)*, pp. 173–182.

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